# Making all equalities equal

Guilherme Horta Alvares da Silva

May 12, 2022

#### Abstract

# **1** Equalities

There are multiple ways of defining equalities in a theorem prover. In the next sections, they will be defined.

### 1.1 Imports

First, it will be necessary to give some agda arguments:

```
{-# OPTIONS -- cubical -- cumulativity #-} module paper where
```

The cubical flag is necessary because we are using cubical equality, and the cumulative flag is also necessary for level subtyping,

open import Agda.Primitive.Cubical using (I; i0; i1)

Private variables are required, so it is not necessary to redefine them later as an implicit variable.

open import Cubical.Core.Primitives using (Level; *l*-max)

private variable  $\ell \ell'$ : Level A: Set  $\ell$ 

This library loads Cubical Agda Primitives as the equality interval.

### 1.2 Martin-Löf Equality

At the beginning of Agda and in most theorems proves, equality is given by Martin-Löf's definition:

module Martin-Löf { $A : \text{Set } \ell$ } where data  $\_\equiv\_(x : A) : A \rightarrow \text{Set } \ell$  where refl :  $x \equiv x$ 

This equality is very convenient in proof assistances like Agda because it is possible to pattern match using them:

private variable  $x \ y \ z : A$ sym :  $x \equiv y \rightarrow y \equiv x$ sym refl = refl trans :  $x \equiv y \rightarrow y \equiv z \rightarrow x \equiv z$ trans refl refl = refl

But the problem of this equality is that it does not handle extensionality and other axioms very well.

```
module FunExt {A : \text{Set } l} {B : \text{Set } l'} where
open Martin-Löf
funExt-Type = {fg : A \rightarrow B}
\rightarrow ((x : A) \rightarrow fx \equiv gx) \rightarrow f \equiv g
```

### **1.3 Cubical Equality**

To solve this problem, Agda adopted cubical type theory that equality is a function from the path to type:

module CubicalEquality { $A : \text{Set } \ell$ } where postulate PathP : ( $A : I \rightarrow \text{Set } \ell$ )  $\rightarrow A \text{ i0} \rightarrow A \text{ i1} \rightarrow \text{Set } \ell$  $\_=\_ : A \rightarrow A \rightarrow \text{Set } \ell$  $\_=\_ = \text{PathP } \lambda \_ \rightarrow A$ 

From this equality, I will define reflection, symmetry and extensionality:

```
module CubicalResults {A : Set \ell} {B : Set \ell'} where
open import Cubical.Core.Primitives
private variable
x y z : A
refl : x \equiv x
refl {x = x} = \lambda \rightarrow x
```

sym :  $x \equiv y \rightarrow y \equiv x$ sym  $p \ i = p \ (\tilde{i})$ funExt :  $\{f g : A \rightarrow B\}$   $\rightarrow ((x : A) \rightarrow f x \equiv g \ x) \rightarrow f \equiv g$ funExt  $p \ i \ x = p \ x \ i$ 

The operator ~ invert the interval. If the interval *i* goes from *i0* to *i1*, the interval ~ *i* goes from *i1* to *i0*.

#### 1.4 Leibniz equality

Leibniz equality is defined in this way: If a is equal to b, then for every propositional P, if P a, then P b. The main idea is that if both values are equal, then they are seen equal for every angle.

```
module LeibnizEquality {A : Set} where

\_\doteq\_: A \to A \to Set_1

a \doteq b = (P : A \to Set) \to P \ a \to P \ b
```

# 2 Joining all equalities

All equalities have something in common. They are all equal to each other. So it will be defined as a common record that all equalities should have. In the next definition, all equalities are equal to cubical equality:

```
open import Cubical.Foundations.Prelude
open import Cubical.Foundations.Isomorphism
open import Cubical.Foundations.Equiv
open import Cubical.Foundations.Univalence
open import Cubical.Foundations.Function
open import Cubical.Data.Equality
module _{-} {\ell} {A : Set \ell'} where
  \triangleq-Type = A \rightarrow A \rightarrow Set \ell
  private
    \ell_1 = \ell-max \ell' \ell
  private variable
    x y z : A
  record IsEquality (\_=: =-Type) : Set (\ell-suc \ell_1) where
    constructor eq
    field
       ≜-≡-≡ : let
```

```
x \equiv y: Type \ell_1
            x \equiv y = x \equiv y
            x \triangleq y: Type \ell_1
            x \triangleq y = x \triangleq y
            in \_\equiv_{} \{\ell \text{-suc } \ell_1\} x \triangleq y x \equiv y
   ≡-≡-≜ : let
         x \equiv y: Type \ell_1
         x \equiv y = x \equiv y
         x \triangleq y = x \triangleq y
         in x \equiv y \equiv x \triangleq y
   ≡-≡-≜ = sym ≜-≡-≡
module _{\{\_} \triangleq_{-} : \triangleq-Type} where
   sym-Equality : (\equiv = = = : \{x \ y : A\} \rightarrow  let
        x \equiv y: Type \ell_1
         x \equiv y = x \equiv y
         x \triangleq y = x \triangleq y
         in x \equiv y \equiv x \triangleq y)
         \rightarrow IsEquality _\triangleq_
   sym-Equality \equiv -\equiv -\triangleq = eq (sym \equiv -\equiv -\triangleq)
record Equality : Set (\ell-suc \ell_1) where
   constructor eqC
   field
         _≜_ : ≜-Type
         { isEquality } : IsEquality ____
EqFromInstance : \{ \triangleq : \triangleq-Type\} \rightarrow IsEquality \triangleq \rightarrow Equality
EqFromInstance inst = eqC \{ inst \}
eqsEqual : (\_\triangleq_{1-}\_\triangleq_{2-}: \triangleq-Type)
   \{ \triangleq_1 - eq : \mathsf{IsEquality}_{\_} \triangleq_{1-} \}
   \{ \triangleq 2 - eq : \mathsf{IsEquality} \_ \triangleq 2 - \}
   \rightarrow \forall \{x \ y\} \rightarrow \mathsf{let}
        x \triangleq _1 y : Type \ell_1
        x \triangleq y = x \triangleq y
         x \triangleq _2 y : Type \ell_1
         x \triangleq _2 y = x \triangleq _2 y
         in \_\equiv_{} \{\ell \text{-suc } \ell_1\} x \triangleq _1 y x \triangleq _2 y
eqsEqual _ _ { eq \triangleq = = = 1 } { eq \triangleq = = = 2 } = \triangleq = = 1 • sym \triangleq = = 2
```

It will be defined for each equality, its instance:

#### 2.1 Cubical Equality

The simplest example is cubical equality hence this equality is already equal in itself.

```
module _ {a} {A : Set a} where

instance

\equiv-IsEquality : IsEquality {A = A} \_\equiv\_

\equiv-IsEquality = eq refl

\equiv-Equality : Equality {\ell = a}

\equiv-Equality = eqC \_\equiv\_
```

#### 2.2 Martin-Löf equality

The proof of Martin-Löf equality is more difficult, but it is already in the Cubical library as p-c.

```
instance

\equiv p-IsEquality : IsEquality {A = A} \_\equiv p_-

\equiv p-IsEquality = sym-Equality p-c

\equiv p-Equality : Equality {\ell = a}

\equiv p-Equality = eqC \_\equiv p_-
```

#### 2.3 Isomorphism

Isomorphism is equality between types.

```
\begin{array}{l} \mbox{module}_{-}\left\{\ell\right\} \mbox{where} \\ \mbox{univalencePath}': \left\{A \ B : \mbox{Type} \ \ell\right\} \rightarrow (A \equiv B) \equiv (A \simeq B) \\ \mbox{univalencePath}' \left\{A\right\} \left\{B\right\} = \\ \mbox{ua} \left\{\ell \mbox{-suc} \ \ell\right\} \left\{A \equiv B\right\} \left\{A \simeq B\right\} (\mbox{compEquiv} (\mbox{univalence} \left\{\ell\right\} \left\{A\right\} \left\{B\right\}) \\ \mbox{(isoToEquiv} (\mbox{iso} \left\{\ell\right\} \left\{\ell \mbox{-suc} \ \ell\right\} \\ \mbox{($\lambda$ $x \to $x$)} (\mbox{$\lambda$ $x \to $x$)} (\mbox{$\lambda$ $b$ $i \to $b$)} \mbox{$\lambda$ $a$ $i \to $a$)})) \end{array}
```

univalencePath is already defined in Agda library, but with  $A \simeq B$  instead of *Lifted*  $(A \simeq B)$ . This change can be done because of the cumulative flag.

instance  $\approx$ -IsEquality : IsEquality  $\{A = \text{Type } \ell\}_{\simeq_{-}}$   $\approx$ -IsEquality = sym-Equality univalencePath'  $\approx$ -Equality : Equality  $\{\ell = \ell\}$  $\approx$ -Equality = eqC\_ $\simeq_{-}$ 

#### 2.4 Leibniz Equality

The hardest equality to prove that is equality is the Leibniz Equality.

```
liftlso : {A : Type \ell} {B : Type \ell'}

\rightarrow lso {\ell} {\ell'} A B \rightarrow lso {\ell-max \ell \ell'} {\ell-max \ell \ell'} A B

liftlso f = iso fun inv

(\lambda x i \rightarrow rightlnv x i) (\lambda x i \rightarrow leftlnv x i)
```

This liftlso will be used to lift the Isomorphism to types of the same maximum level of both.

```
where open Iso f
open import leibniz
open Leibniz
```

It is importing the definition of Leibniz equality made by [?]. In this work, there is already proof of the isomorphism between Leibniz and Martin-Löf equality.

```
module FinalEquality {A : Set} where
open MainResult A
private variable
x y z : A
\doteq \cong \equiv : Iso (x \doteq y) (x \equiv p y)
\doteq \cong \equiv =  iso j i (ptoc \circ ji) (ptoc \circ ij)
```

In Cubical Library, the definition of isomorphism uses cubical equality instead of Martin-Löf equality when we have to prove that  $\forall x \rightarrow from (to x) \equiv x$  and  $\forall x \rightarrow to (from x) \equiv x$ . ptoc is necessary to do this conversion from these equalities.

```
 \stackrel{i}{=}\equiv : (x \stackrel{i}{=} y) \equiv c \ (x \equiv p \ y)  \stackrel{i}{=}\equiv = let \ lifted = liftlso \stackrel{i}{=}\cong in \ isoToPath \ lifted
```

Using the univalence and liftlso defined previously, it is possible to transform the isomorphism into equality.

```
open IsEquality

instance

\doteq-IsEquality : IsEquality \{A = A\}_{==}^{=}_{=}_{=}^{=}_{=}

\doteq-IsEquality = eq \lambda \{x\} \{y\} \rightarrow \pm \equiv \equiv \bullet

\lambda i \rightarrow \equiv p-IsEquality \{\ell \text{-zero}\}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{=}^{=}_{}
```

The last pass is to join the three equalities between equalities: Leibniz to Martin-Löf to cubical equality.

### **3** New Equalities types

The equalities used previously were defined using cubical equality. Now I will define them using other equalities.

```
module Equalities \{a \ \ell\} \ \{A : \text{Set } a\} where
private
\triangleq-Type' = \triangleq-Type \{a\} \ \{\ell\} \ \{A\}
\ell_1 = \ell-max a \ \ell
private variable
x \ y \ z : A
```

Loaded the modules using the levels to be more generic.

```
\begin{array}{l} \text{module} \_ \\ (Eq_1 : \text{Equality} \{\_\} \{\ell\} \{A\}) \\ \text{where} \end{array}
```

```
open Equality Eq 1 renaming (\_=\_ to \_= 1-; is Equality to eq 1)
```

I am importing generic equality to use it to define more generic equality.

```
record IsEquality _2 (\_=\_:=-Type') : Set (\ell-suc \ell_1) where
constructor eq
field
\triangleq_{====:} : let
x \equiv y : Type \ell_1
x \equiv y = x \equiv _1 y
```

Different from previously definition of IsEquality, the cubical equality defined in the line above was substituted by the more generic equality  $\equiv 1$ .

```
x \triangleq y : \text{Type } \ell_1

x \triangleq y = x \triangleq y

in _=_ {\ell-suc \ell_1} x \triangleq y x \equiv y

=-==-\triangleq : let

x \equiv y : \text{Type } \ell_1

x \equiv y = x \equiv 1 y

x \triangleq y = x \triangleq y

in x \equiv y \equiv x \triangleq y

\equiv -\equiv -\triangleq = \text{sym} \triangleq -\equiv -\equiv
```

The rest of the definition is the same.

```
instance
≜-isEquality : IsEquality _≜_
≜-isEquality = eq (≜-≡-≡ ● IsEquality.≜-≡-≡ eq 1)
```

From a more generic definition of equality, it is easily possible to return to the less generic definition.

```
module _

(Eq_2 : Equality {_} {\ell} {\ell} {A})

where

open Equality Eq_2 renaming (_\triangleq_ to _= 2_; isEquality to eq 2)
```

I am defining a new generic equality to prove that it is an equality of type 2:

```
eqsEqual 2 : let
   x \triangleq _1 y: Type \ell_1
  x \triangleq _1 y = x \equiv _1 y
   in x \triangleq y \equiv x \equiv y
eqsEqual_2 = eqsEqual_{= 1-} = 2-
instance
   \equiv 2-Equality 2 : IsEquality 2 \_\equiv 2-
   \equiv 2-Equality 2 = eq (sym eqsEqual 2)
      where open IsEquality
module \_ \{\_ \triangleq \_ : \triangleq -Type\} where
   sym-Equality 2 : (\equiv -\equiv -\triangleq : \{x \ y : A\} \rightarrow let
      x \equiv y: Type \ell_1
      x \equiv y = x \equiv 1 y
      in x \equiv y \equiv (x \triangleq y)
      \rightarrow IsEquality 2 _4_
   sym-Equality _2 \equiv -\equiv -\triangleq = eq (sym \equiv -\equiv -\triangleq)
```

Given a symmetric definition of the previous equality, it is easy to prove that it is also equality of type 2.

#### **3.1** Everything is an equality

In this part, a relation is equality when it is equal (using general equality) to cubical equality.

```
module _
(Eq_3 : Equality {A = \text{Set } \ell_1})
where
```

open Equality Eq 3 renaming ( $\_=\_$  to  $\_=$  3-; isEquality to eq 3)

```
record IsEquality _3 (_\triangleq_- : \triangleq-Type') : Set (\ell-suc \ell_1) where
constructor eq
field
\triangleq-\equiv-\equiv : (x \triangleq y) \equiv_3 (x \equiv_1 y)
```

With this definition of equality, it is possible to prove that if equality is equal to cubical equality, so it is equal (using the general or cubical equality) to the cubical equality.

This is proof that the symmetric definition of equality is also valid.

```
 = -= - \triangleq : (x \equiv _1 y) \equiv _3 (x \triangleq y) 
 = -= - \triangleq = let 
 \alpha_1 = lsEquality. = -= - \triangleq eq_3 
 \alpha_2 = lsEquality. = -= = eq_1 
 \alpha_3 = lsEquality. = -= - \triangleq = -isEquality 
in transport \alpha_1 (\alpha_2 \bullet \alpha_3)
```

It is possible to prove that a general equality is equality from this definition:

```
module _

(Eq_2 : Equality \{.\} \{.\}\}) where

open Equality Eq_2 renaming (\_=\_ to \_= \_\_]; isEquality to eq_2)

instance

\equiv \__2-Equality \__3 : IsEquality \__3 \_= \__2-

\equiv \__2-Equality \__3 = eq \alpha

where

open IsEquality eq_3

\alpha : (x \equiv \__2 y) \equiv \__3 (x \equiv \__1 y)

\alpha = transport \equiv -\equiv -\triangleq (IsEquality \__2.\triangleq -\equiv -\equiv (\equiv \_2-Equality \__2 (eqC \_= \__2)))
```

If there is proof of symmetrical equality, so it is also equality from this definition:

module  $_{-} \{\_ \triangleq_{-} : \triangleq_{-} Type'\}$  where

```
sym-Equality <sub>3</sub> :

(\equiv -\equiv -\triangleq : \forall \{x \ y\} \rightarrow (x \equiv_1 y) \equiv_3 (x \triangleq y))

\rightarrow IsEquality <sub>3</sub> \_=\_.

sym-Equality <sub>3</sub> \equiv -\equiv -\triangleq = eq (let

\alpha_1 = IsEquality.\equiv -\equiv -\triangleq eq_3

\alpha_2 = transport (sym \alpha_1) \equiv -\equiv -\triangleq

in transport \alpha_1 (sym \alpha_2))
```

## **4** Using the definitions

The best part of defining all of this stuff is that it is now easy to prove that Leibniz equality is equality.

```
module LeibnizFromPEquality {A : Set} where
open Equalities {\ell-zero} {\ell-suc \ell-zero}
\_\equiv p_{1-}: A \rightarrow A \rightarrow Set_1
x \equiv p_1 y = x \equiv p y
```

I redefined this equality because it must be a set of universe one. And because of that, I have to prove again that this is an equality:

```
instance

\equiv p_{1}\text{-isEquality} : \text{IsEquality}_{-}\equiv p_{1-}
\equiv p_{1}\text{-isEquality} = \text{eq } \lambda \{x \ y\} \rightarrow (\text{sym } \lambda \ i \rightarrow \text{let}
\alpha : \text{Type}_{1}
\alpha = p\text{-c} \{\ell\text{-zero}\} \{x = x\} \{y = y\} i
in \alpha)
```

With just one line of code, it is possible now to prove that Leibniz equality is an equality from Martin-Löf Equality.

leibniz : IsEquality  $\{A = A\}_{=}^{=}$ leibniz = IsEquality  $_{2}$ .=-isEquality  $\{Eq_{1} = eqC_{=}p_{-}\}$  (eq FinalEquality.==)